AN IMPERATIVE REPETITION FOR VEHICLE AD-HOC NETWORK COMMUNICATION SYSTEM

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ABSTRACT

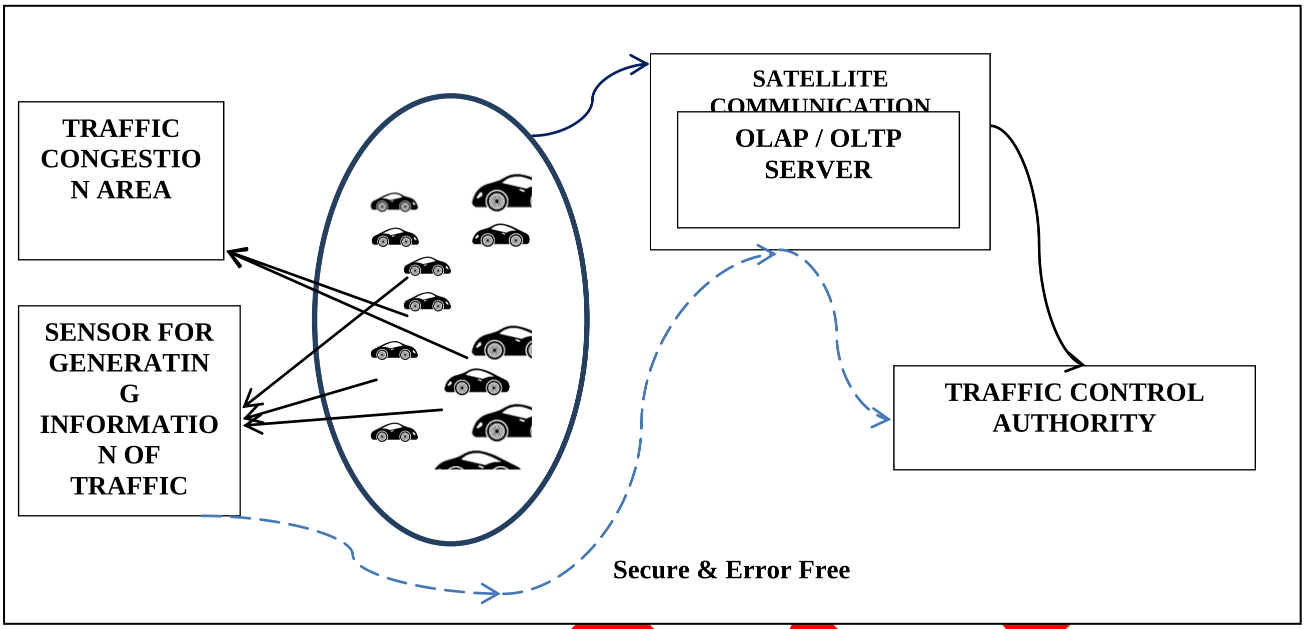
QoS of VANET is too much tedious task for humans, as life challenging. In this article we approach a system for prevent congestion as well as recovery after congestion. Today's era without communication, the life is almost not possible to continue in form of luxuries and satisfaction, we introduce a new approach for secure and error free communication for congestion free VANET.

Index Terms—VANET, ORDES, Fuzzy Error Correction Code, Communication Security.

I. INTRODUCTION

This is a system that will help out all the persons in form of congestion free VANET. We introduce a system for VANET using various current technology. Basically it is an architecture for implementation using communication network MANET, Mobile Agent and Sensors.

In this architecture the system follow all the basic aspects of communication system and also communicate with our helping hands like (Traffic Department, Fire Department and congestion Recovery etc.). The flow of information is necessary to know the system Ex. Information for Recovery agencies as could as various departments for downwards and sending report to higher authority as upward flow for making efficient decision to instruct the agencies. In this architecture we analyses the congestion over traffic on roads. The system always gives the effective solution and provide help for person, thing resource which are located at congestion Site. Basically the system is a VANET i.e. Vehicle Ad-Hoc Network using high effective sensors with connecting sub networks like Personal Area Network (PAN), Storage Area Network (SAN). Using SAN we create a server which will work on OLAP, OLTP. The server consists of all information for Disaster and have some authority to generate effective decision with help of Artificial Intelligence expert systems.



The complete system are divided into three parts:

1. Congestion Site
2. VANETAuthority
3. Communication Channel (Secure and Error Less)

II. PRELIMINARIES

2.1 VANET

VANETs are distributed, self-organizing link webs crafted up from voyaging vehicles, and are consequently delineated by tremendously elevated speed and manipulated degrees of freedom in nodes movement patterns. Such particular features oftentimes make average networking protocols inefficient or unusable in VANETs, and this, joined alongside the huge encounter that the arrangement of VANET technologies could have on the automotive marketplace, explains the producing manipulation in the progress of link protocols that are specific to vehicular networks. The frank believed of VANET is straightforward: seize the extensively adopted and inexpensive wireless innate span web (WLAN) knowledge that links notebook computers to every single supplementary and the Internet, and, alongside an insufficient tweaks, installed on the vehicles. Of sequence, if it were honestly that unambiguous, the alert. VANET scrutiny area should probable not ever have formed. Vehicular environment creates exceptional opportunities ,trials, and requirements. If vehicles can undeviatingly converse alongside every single supplementary and alongside groundwork, a jointly new prototype for vehicle protection requests can be created. Even supplementary non-safety requests can rise road and vehicle efficiency. Second, new trials are crafted by elevated vehicle speeds and exceedingly vibrant working environments. Third ,new necessities, essential ised by new safety-of-life proposition, contain new outlook for elevated packet transfer rates and low packet latency. Further, client agreement and governmental oversight hold extremely elevated expectations of privacy and security. Even nowadays, vehicles produce and examine colossal numbers of data, even though normally this data is self-collected inside a solitary vehicle. With a VANET, the 'horizon of awareness' for the vehicle or driver drastically grows. The VANET contact can be which ever completed undeviatingly amid vehicles as 'one-hop' contact, or vehicles can retransmit memos, thereby enabling 'multihop' communication. To rise coverage or robustness of contact, relays at the roadside can be deployed. Roadside groundwork can additionally be utilized as a gateway to the Internet and, therefore, data and context data can be amassed, stored and processed' somewhere', e.g., in Cloud infrastructures. The earth of vehicular request and inter-networking technologies is established on an interdisciplinary power in the cross serving of contact and networking, automotive electronics, road procedure and association, and data and ability provisioning. VANET can consequently be perceived as an vital portion of intelligent transportation arrangements(ITS). Vehicular Ad-Hoc Web (VANET) contact has presently come to be an increasingly accepted scrutiny case in the span of wireless networking as well as the automotive industries. The aim of VANET scrutiny is to develop a vehicular contact arrangement to enable quick and cost-efficient allocation of data for the benefit of passengers' protection and comfort. VANETs need specific networking methods alongside feasibility and performance.

2.2 ORDES Approach:

We also using the same feistel structure and same process for encryption and decryption but we add a new process for key generation. In this process, key itself generate n different keys using a function and random number generated by Hardware Random Number Generator (HRNG) then new generated key block applies on the each block of message for all round of DES. For each block of message, the process generates a separate key. This new generated key used in encryption phase as well as the decryption phase.

M = {m1, m2, m3, ..., mn} (64- bit block of message)
K = {K} (56- bit Key)

2.2.1Key Generation:

F {K and Rj} = [K_new (x)_p]

Where Rj generated by Hardware Random Number Generator (HRNG) and [1 ≤ Rj ≤ (256 = 72,057,594,037,927,936)].

[K_new (x)_p)] = [K_new (x)_1] to [K_new (x)_n]

Function F:

Step 1- Input the bit value of initial key K (56-bit).
Step 2- Input generated random number Rj , generated by HRNG*.
(*HRNG Property- 256 no., random number generator)
Step 3- Convert Rj into 56- bit binary number.
Step 4- Now, we have
Key K = {KB1, KB2, KB3, ..., KB56}
And Rj ={Rb1, Rb2, Rb3, ..., Rb56}
Where KBr is the bit of Key and Rbr is the bit of Random no. Here r =1, 2, 3...............56.
Step 5- Apply condition on K and Rj.
IF Rbr = 1 then, Complement (convert 1 to 1 or 0 to 0) of corresponding KBr.
ANDIF Rbr = 0 then, Retain the same (1 to 1 or 0 to 0) of corresponding KBr.
Step 6- [K_new (x)_p] = Result of step 5.
Using this function F every time we get the result [K_new (x)_p] for each block of message. For each block of M we generate a new no. Rj and implement function F. Finally get a new key for every block of message.

2.2.2 Encryption/Decryption:

In encryption phase, ORDES take a message block mn and a new generated key [K_new (x)_p], implement encryption process as per traditional DES. One special thing make aur process is different- REKEYING PROCESS.
REKEYING PROCESS have the property to make various key for various block of message.
Now, we have a new key for every block of message. This new key [K_new (x)_p] is apply on each block of message M.
In this process, new key is also make 16 different key for every round of DES using shifting property as per traditional DES. For every block of message M, new key [K_new (x)_p] makes a new key block for every round of DES to implement in the encryption process.

Decryption Process is the inverse step of encryption process. In decryption, we also use the same key which is used in encryption.
Ci = E_([K_new (x)_p)] {mi} and mi = D_([K_new (x)_p]) {Ci}, where $1 \leq i \leq n$.
Cipher Text C = {C1, C2, ..., Cn} and
Plain Text M = {m1, m2, ..., mn}.

2.3. Crisp Commitment Schemes

In a commitment scheme, one party sender aim to entrust a concealed message m to the second party receiver, intuitively a commitment scheme may be seen as the digital equivalent of a sealed envelope. If Sender wants to commit to some message m, sender just puts it into the sealed envelope, so that whenever sender wants to reveal the message to receiver, sender opens the envelope. First of all the digital envelope should hide the message from: receiver should be able to learn m from the commitment. Second, the digital envelope should be binding, meaning with this that sender can not change her mind about m, and by checking the opening of the commitment one can verify that the obtained value is actually the one sender had in mind originally.

2.3.1 Definition: A Commitment scheme is a tuple$\{P,E,M\}$ Where $M=\{0,1\}^n$ is a message space, P is a set of individuals , generally with three elements A as the committing party, B as the party to which Commitment is made and TC as the trusted party , $E=\{(t_i,e_i)\}$ are called the events occurring at times t_i , $i=1,2,3$, as per algorithms e_i , $i=1,2,3$. The scheme always culminates in either acceptance or rejection by A and B.

The environment is setup initially, according to the algorithm$setupalg\ (e_1)$ and published to the parties A and B at timet_1.

During the Commit phase, A uses algorithm$commitalg\ (e_2)$, which encapsulates a message$m\in M$, along with secret string $S\in_R \{0,1\}^k$ into a string$c$. The opening key (secret key) could be formed using both m and$S$. A sends the result c to B (at timet_2).

In the Open phase, A sends the procedure for revealing the hidden Commitment at time t_3, and B uses this. $openalg\ (e_3)$: B constructs c'using $commitalg$, message m and opening key, and checks weather the result is same as the commitment c .

Decision making:

If $c=c'$.

Then A is bound to act as in m

Else he is free to not act as m

2.3.2 Definition: Fuzzy Commitment scheme is a tuple$\{P,E,M,f\}$ Where $M\subseteq\{0,1\}^k$ is a message space which consider as a code, P is a set of individuals , generally with three elements A as the committing party, B as the party to which Commitment is made and TC as the trusted party , f is error correction function (def. 2.2.5) and $E=\{(t_i,e_i)\}$ are called the events occurring at times t_i , $i=1,2,3$, as per algorithms e_i , $i=1,2,3$. The scheme always culminates in either acceptance or rejection by A and B.

In the setup phase, the environment is setup initially and public commitment key CK generated, according to the algorithm$setupalg\ (e_1)$ and published to the parties A and B at timet_1.

During the Commit phase, Alice commits to a message $m\in M$ according to the algorithm$commitalg\ (e_2)$ into string$c$.

In the Open phase, A sends the procedure for revealing the hidden Commitment at time algorithmt_3 and B uses this algorithm$openalg\ (e_3)$: B constructs c' using algorithm$commitalg$, message t (m) and opening key, and checks weather the result is same as the received commitment$t(c)$, where t is the transmission function. Fuzzy decision making:

If $(nearest(t(c),f(c'))\le z_0$

Then A is bound to act as in m

Else he is free to not act as m

2.3.3 Definition: A metric space is a set C with a distance function dist : $C\times C\to R^+=[0,\infty)$, which obeys the usual properties(symmetric, triangle inequalities, zero distance between equal points).

2.3.4 Definition: Let $C\{0,1\}^n$ be a code set which consists of a set of code words c_i of length n. The distance metric between any two code words c_i and c_j in C is defined by

$$dist(c_i, c_j) = \sum_{r=1}^{n} |c_{ir} - c_{jr}| \qquad c_i, c_j \in C$$

This is known as Hamming distance.

2.3.5 Definition: An error correction function f for a code C is defined as $f(c_i) = \{c_j \,/\, dist(c_i, c_j) \text{ is the minimum, over } C - \{c_i\}\}$. Here, $c_j = f\left(c_i\right)$ is called the nearest neighbor of c_i.

2.3.6 Definition: The measurement of nearness between two code words c and c' is defined by $\text{nearness } (c, c') = dist(c, c')/n$, it is obvious that $0 \le \text{ nearness } (c, c') \le 1$.

2.3.7 Definition: The fuzzy membership function for a codeword c' to be equal to a given c is defined as

$$\begin{aligned} FUZZ(c') &= 0 \qquad && \text{if nearness}(c, c') = z \le z_0 < 1 \\ &= z && \text{otherwise} \end{aligned}$$

2.4 SEQUITUR Algorithm

The SEQUITUR algorithm represents a finite sequence _ as a context free grammar whose language is the singleton set {σ}. It reads symbols one-by-one from the input sequence and restructures the rules of the grammar to maintain the following invariants:

(A) no pair of adjacent symbols appear more than once in the grammar, and

(B) every rule (except the rule defining the start symbol) is used more than once. To intuitively understand the algorithm, we briefly describe how it works on a sequence 123123. As usual, we use capital letters to denote non-terminal symbols. After reading the first four symbols of the sequence 123123, the grammar consists of the single production rule S ➔ 1, 2, 3, 1 where S is the start symbol. On reading the fifth symbol, it becomes S ➔ 1, 2, 3, 1, 2 Since the adjacent symbols 1, 2 appear twice in this rule (violating the first invariant), SEQUITUR introduces a non-terminal A to get

S ➔ A, 3,A A ➔1, 2

Note that here the rule defining non-terminal A is used twice. Finally, on reading the last symbol of the sequence 123123 the above grammar becomes

S ➔ A, 3, A, 3 A ➔ 1, 2

This grammar needs to be restructured since the symbols A, 3 appear twice. SEQUITUR introduces another non-terminal to solve the problem. We get the rules

S → B,B
B → A 3
A→ 1 2

However, now the rule defining non-terminal A is used only once. So, this rule is eliminated to produce the final result.

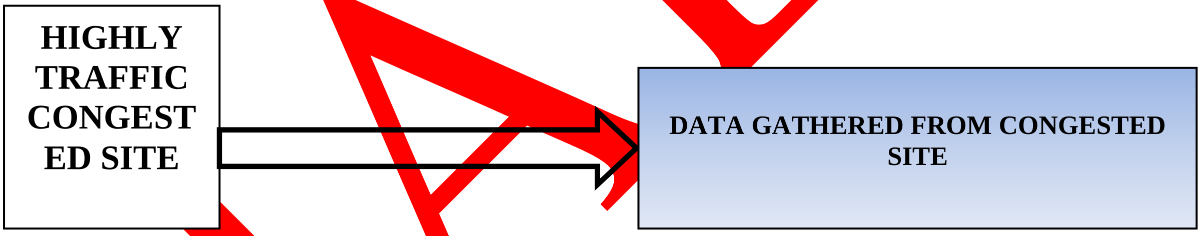
S → B, B B → 1, 2, 3

Note that the above grammar accepts only the sequence 123123.

III. IMPLEMENTATION OF PROPOSED SYSTEM

The implementation of proposed system is divided into these parts.

2.2 Congestion Site: First one establishment of site where are disaster may be occur or in other words they are which are identical for Congestiion occurrence.


HIGHLY TRAFFIC CONGEST ED SITE
DATA GATHERED FROM CONGESTED SITE

Generate data through sensors and transfer data via network may be MANET.

3.2 Secure Communication:

For secure communication we implement ORDES approach with Fuzzy Error Correction Code Now Consider Secure channel for secure communication and the main aspect in reference of unnatural disaster we use fuzzy error correction technique with ORDES Algorithm.

Now apply ORDES Algorithm on Generated Data for High Level Security.
In encryption phase, ORDES take output of compression phase as a message block M_n and a new generated key K_(new n) implement encryption process as per traditional DES.
In this process, New key is also make 16 different key for every round of ORDES using shifting property as per traditional DES. For every block of message M, new key K_(new n) makes a new key block for every round of DES to implement in the encryption process.

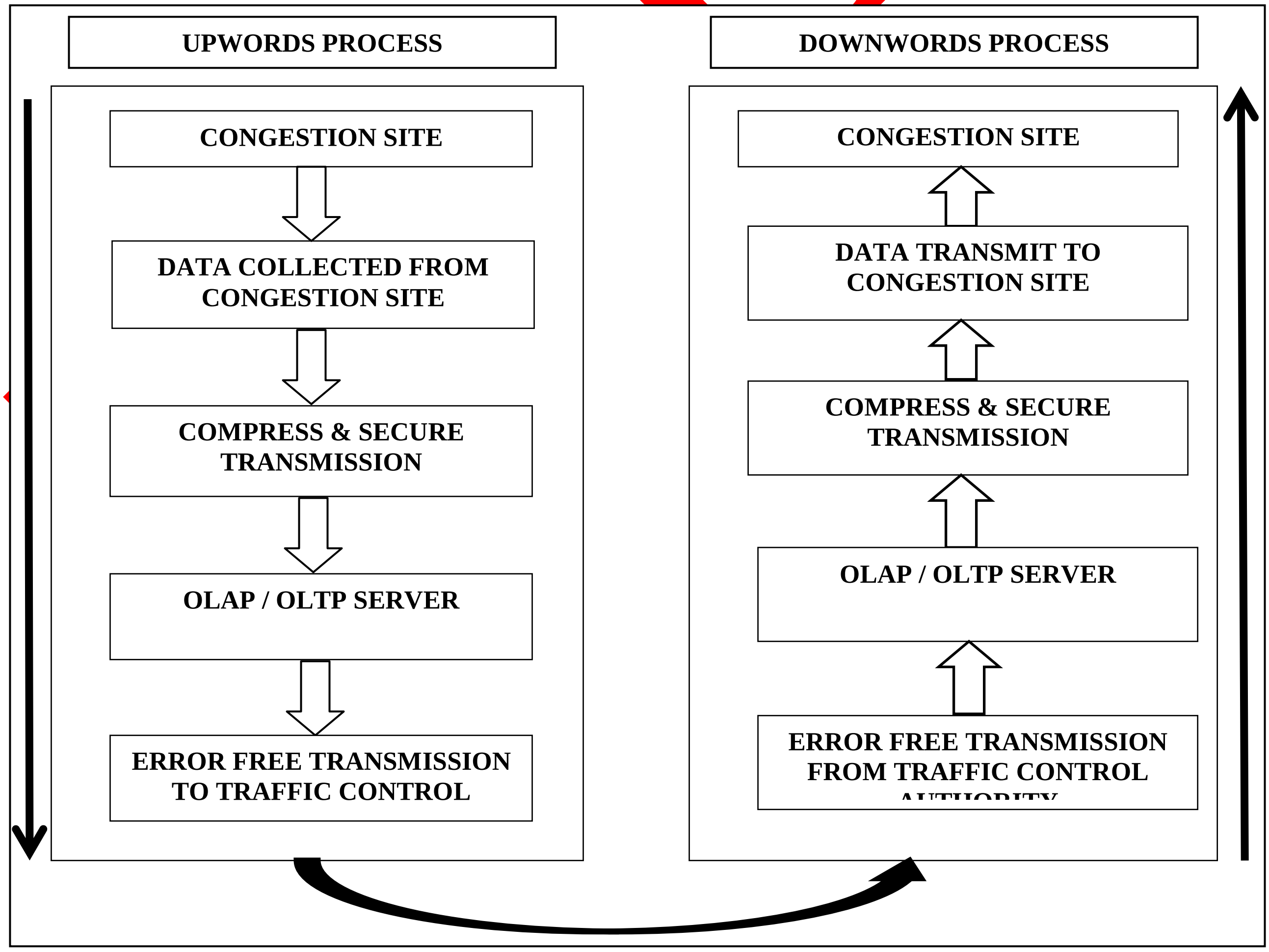
Decryption Process is the inverse step of encryption process. In decryption, we also use the same key which is used in encryption.

〖Ci=E〗_(K_(new n)) {mi} and 〖mi=D〗_(K_(new n))) {Ci},
where $1 \leq i \leq n$.
Cipher Text C={C1,C2………Ci} and Plain Text M={m1,m2………mi}

Compression using SEQUITUR:
After quantization, the scheme uses a filter to pass only the string of non-zero coefficients. By the end of this process we will have a list of non-zero tokens for each block preceded by their count. DCT based image compression using blocks of size 8x8 is considered. After this, the quantization of DCT coefficients of image blocks is carried out. The SEQUITER compression is then applied to the quantized DCT coefficients.

Error Correction:
Receiver check that , he will realize that there is an error occur during the transmission. Receiver apply the error correction function f to .
Then receiver will compute nearness

$$(t(c), f(c')) = dist(t(c)\, f(c'))/n \qquad FUZZ(c') = 0 \quad \text{if } \text{nearness}(c, c') = z \leq z_0 < 1$$
$$= z \quad \text{otherwise}$$


UPWORDS PROCESS
CONGESTION SITE
DATA COLLECTED FROM CONGESTION SITE
COMPRESS & SECURE TRANSMISSION
OLAP / OLTP SERVER
ERROR FREE TRANSMISSION TO TRAFFIC CONTROL
DOWNWORDS PROCESS
CONGESTION SITE
DATA TRANSMIT TO CONGESTION SITE
COMPRESS & SECURE TRANSMISSION
OLAP / OLTP SERVER
ERROR FREE TRANSMISSION FROM TRAFFIC CONTROL AUTHORITY

V. CONCLUSION

In this paper, we implement ORDES scheme, Sequitur and Fuzzy error correction code for Disaster Management. As we know that ORDES provides n-times more security in compare of DES but only in two cases ORDES works like DES.

Case I

If we take one and only one Key on the place of n keys then 〖K_new (x〗 _p) is K, at this condition our approach works like DES.

Case II

If 〖K_new (x〗 _1) = 〖K_new (x〗 _2) = 〖K_new (x〗 _n), than our approach also works like DES.

If, after encryption anyone wants to communicate message with commitment. The encoded message is transmitted. It is possible that during the transmission some bits of are changed. The receiver receives the incorrect message . He solves the decoding problem, that is, he calculates such that is minimum. If the error is not too big, that is, , where is the minimum distance of any two distinct code words, then is equal to the original message .

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